

An introduction to core entropy

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University of Toronto

Summary

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Then the **topological entropy** of f is

$$h_{top}(f) := \sup_{\mathcal{U}} \lim_{n \rightarrow \infty} \frac{1}{n} \log N(\mathcal{U} \vee f^{-1}(\mathcal{U}) \vee \dots \vee f^{-n+1}(\mathcal{U}))$$

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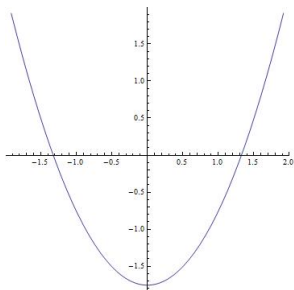
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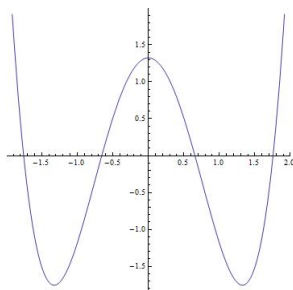
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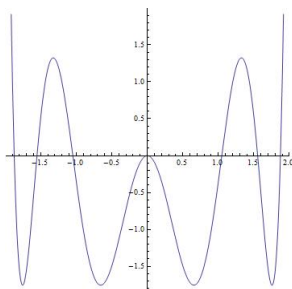
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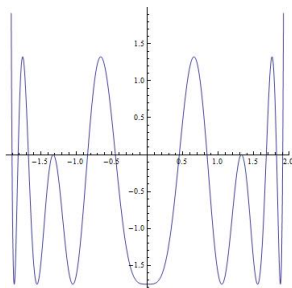
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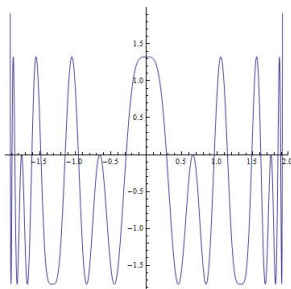
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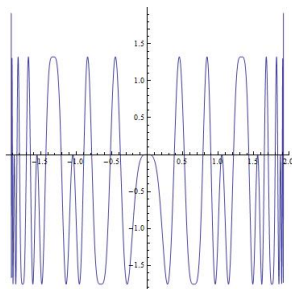
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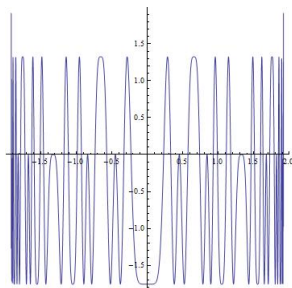
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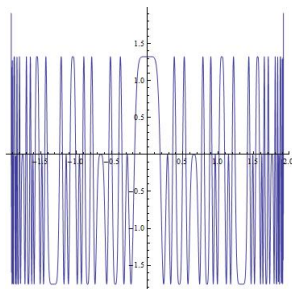
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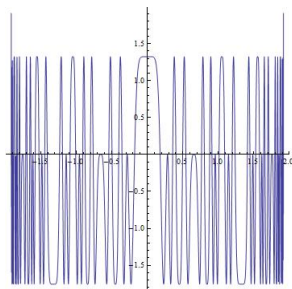
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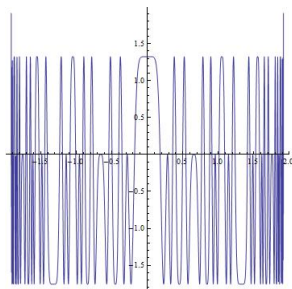
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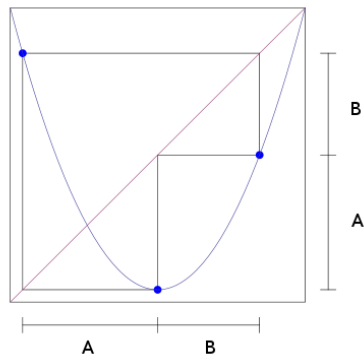


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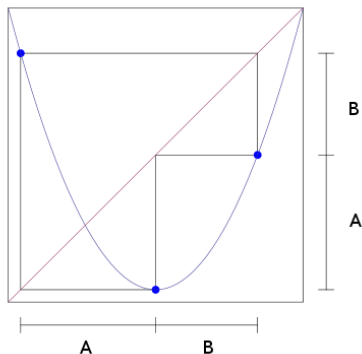
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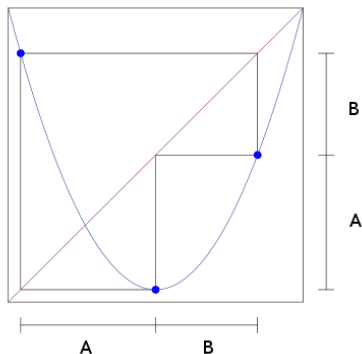


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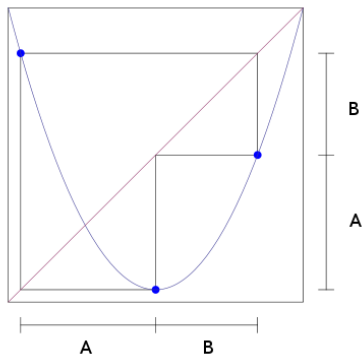
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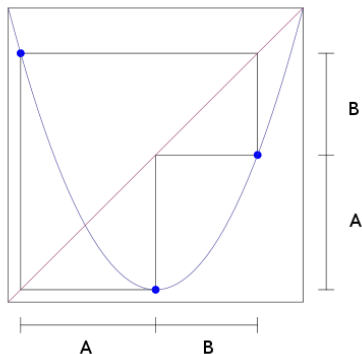
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How does entropy change with the parameter c ?

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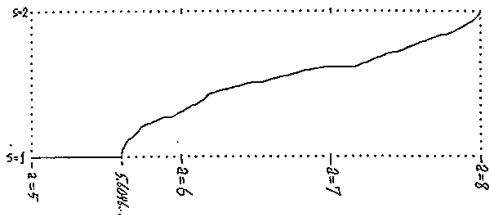
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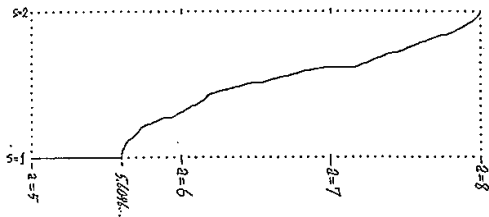
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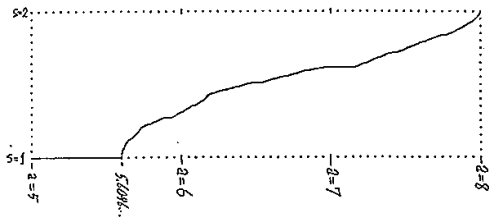
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[Picture is for $f_a(x) = ax(1 - x)$.]

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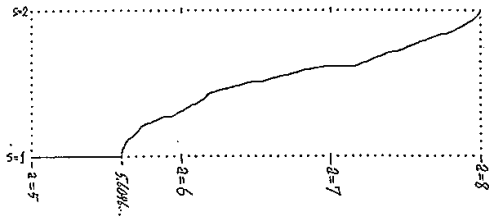
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Question : Can we extend this theory to complex polynomials?

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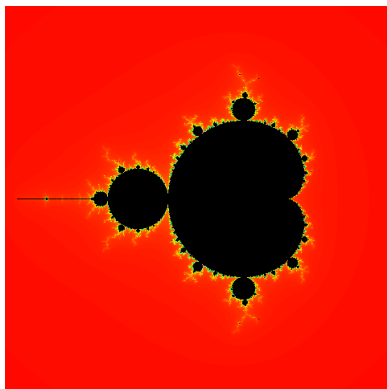


Remark. If we consider $f_c : \hat{\mathbb{C}} \rightarrow \hat{\mathbb{C}}$ entropy is **constant**
 $h_{top}(f_c, \hat{\mathbb{C}}) = \log 2$. (Lyubich 1980)

Mandelbrot set

The **Mandelbrot set** \mathcal{M} is the connectedness locus of the quadratic family

$$\mathcal{M} = \{c \in \mathbb{C} : f_c^n(0) \not\rightarrow \infty\}$$



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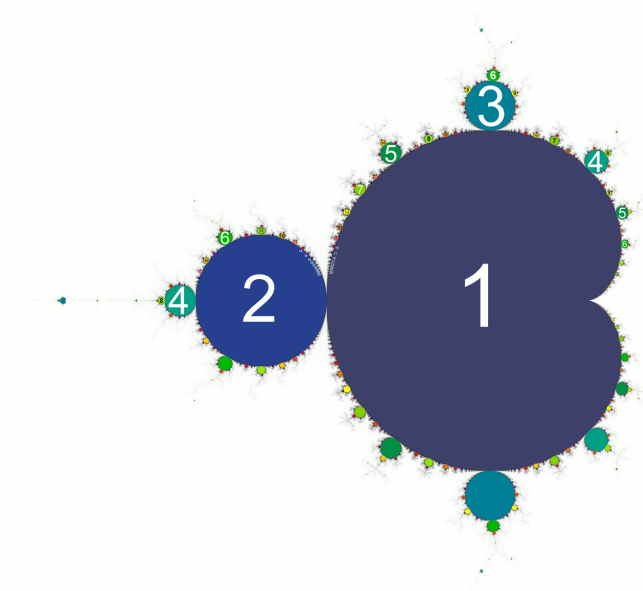
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Each hyperbolic component has a period, and is biholomorphic to the disk.



External rays

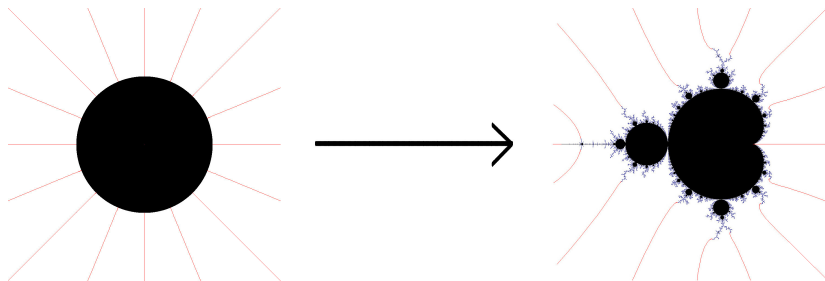
Since $\hat{\mathbb{C}} \setminus \mathcal{M}$ is simply-connected, it can be uniformized by the exterior of the unit disk

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The images of radial arcs in the disk are called **external rays**.
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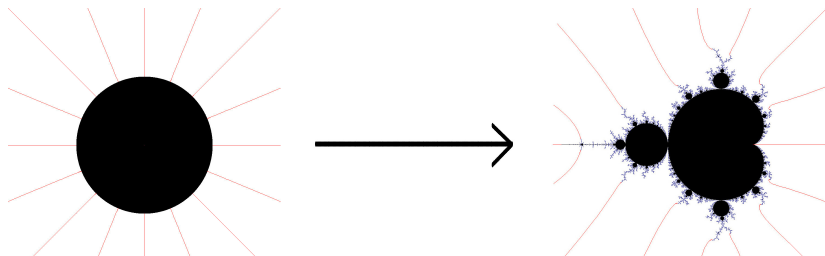
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As a consequence, the Mandelbrot set is homeomorphic to a quotient of the closed disk (hence locally connected).

Julia sets

Let $f_c(z) = z^2 + c$. Then the filled Julia set of f_c is the set of points which do not escape to infinity under forward iteration:

$$K(f_c) := \{z \in \mathbb{C} : f_c^n(z) \text{ is bounded} \}$$

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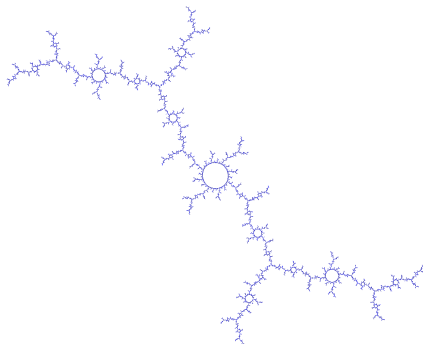
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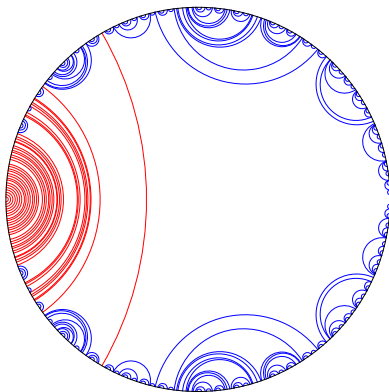
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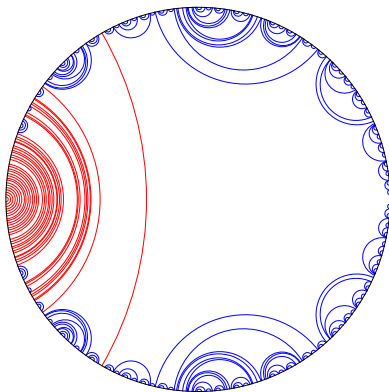
$$f_c(\gamma_c(\theta)) = \gamma_c(2\theta)$$

Thurston's quadratic minor lamination (QML)



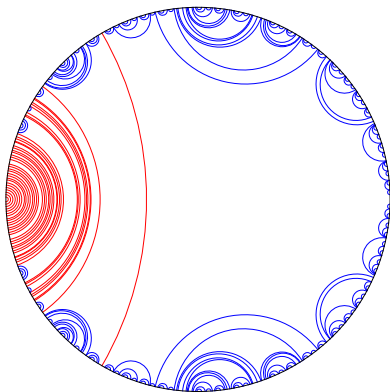
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For each f_c , pick the **minor leaf** of the lamination for f_c (i.e., the ray pair landing at the critical value (or its root)). The **QML** is the union of all minor leaves for all $c \in \mathcal{M}$. The quotient \mathcal{M}_{abs} of the disk by the lamination is a (locally connected) model for the Mandelbrot set, and homeomorphic to it if MLC holds.

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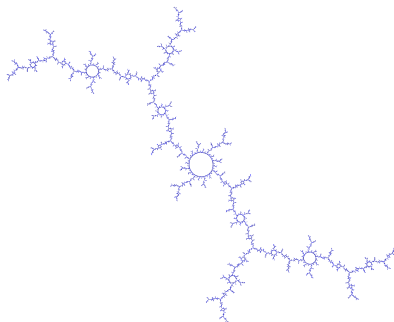
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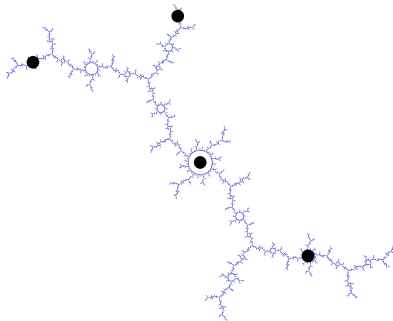


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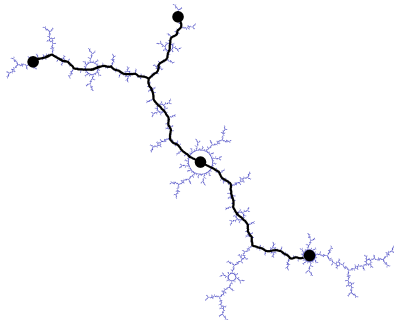


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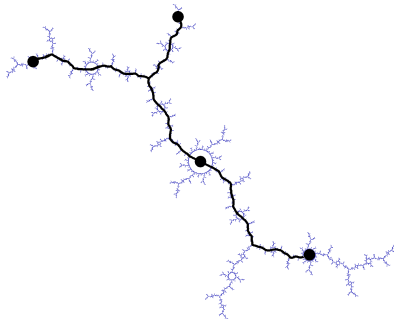


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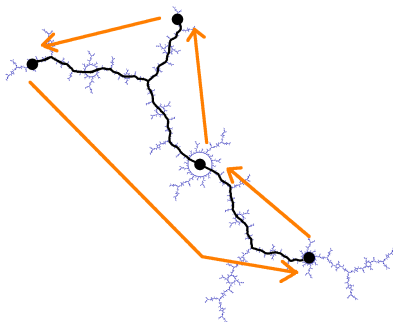


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The core entropy

Definition (W. Thurston)

Let f be a polynomial whose Julia set is connected and locally connected

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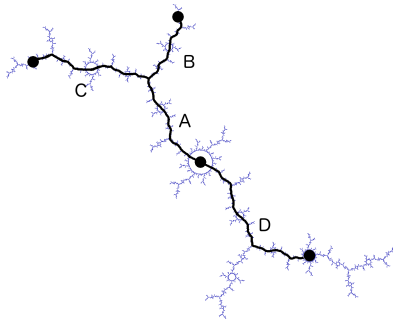
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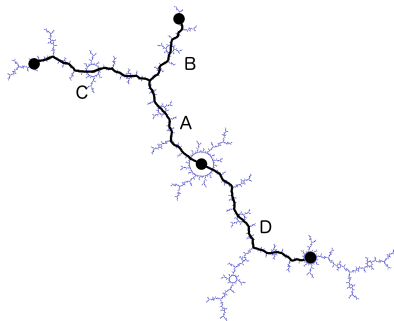
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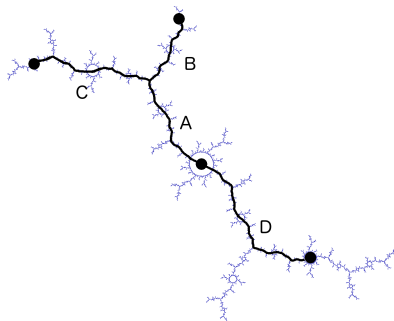
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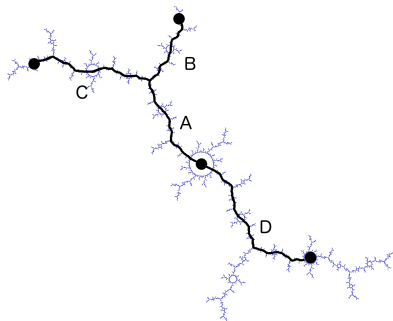


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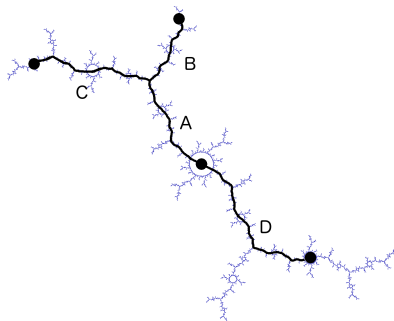
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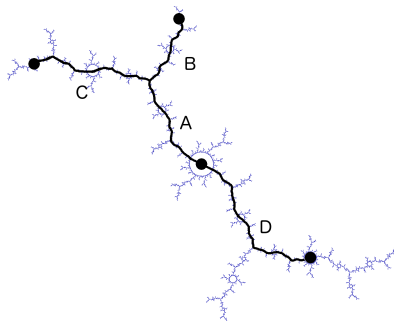
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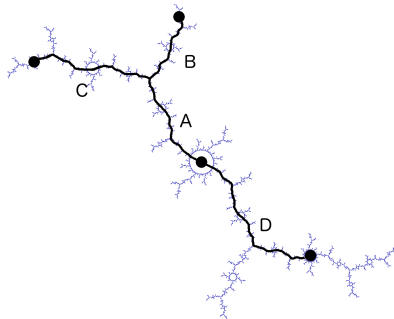


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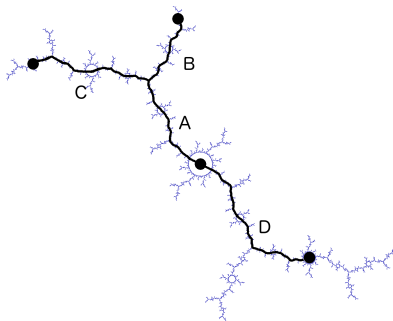
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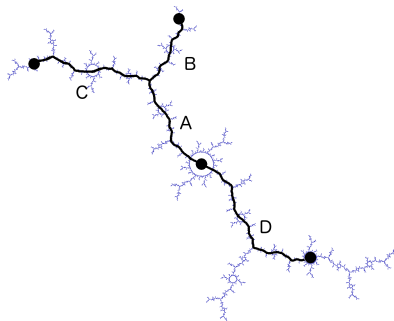
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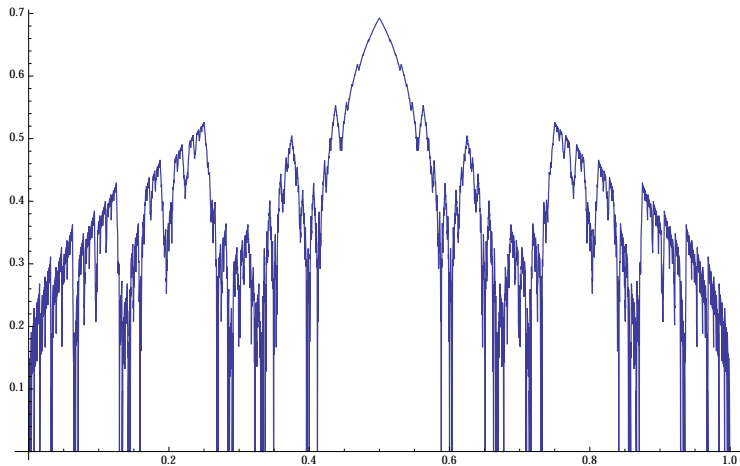
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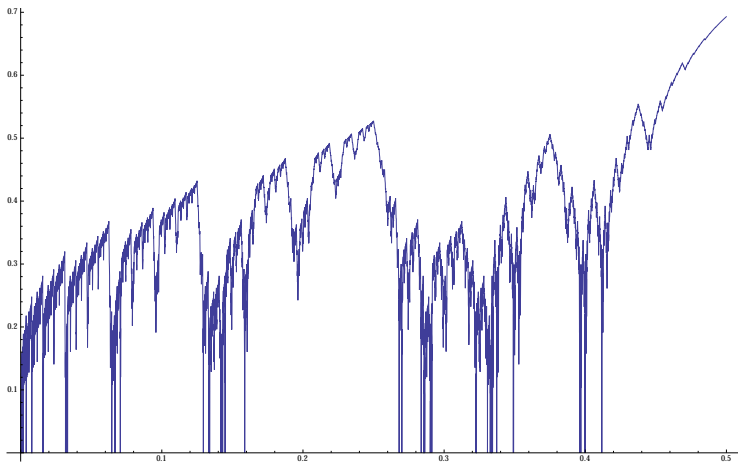
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Question: How does $h(\theta)$ vary with the parameter θ ?

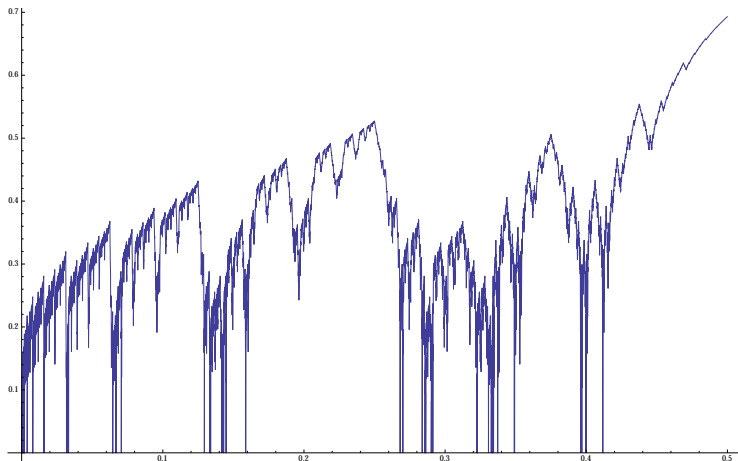
Core entropy as a function of external angle (W. Thurston)



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Question Can you see the Mandelbrot set in this picture?

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Theorem (Li Tao; Penrose; Tan Lei; Zeng Jinsong)

If $\theta_1 <_M \theta_2$, then

$$h(\theta_1) \leq h(\theta_2)$$

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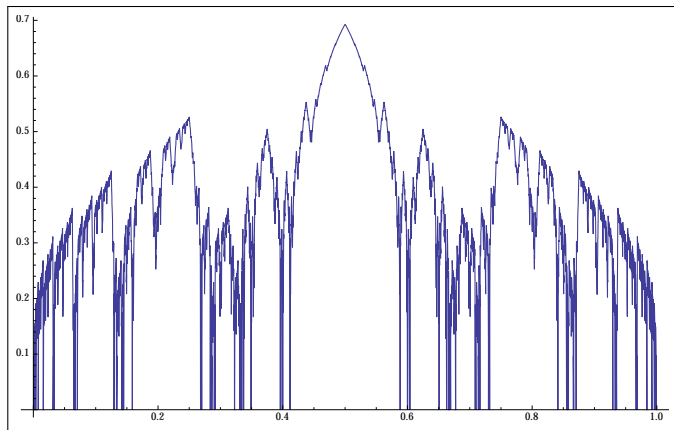
Theorem (T., Bruin-Schleicher)

If the Hubbard tree of f_c is topologically finite, then

$$\text{H. dim } B_c = \frac{h(f_c)}{\log 2}$$

The core entropy as a function of external angle

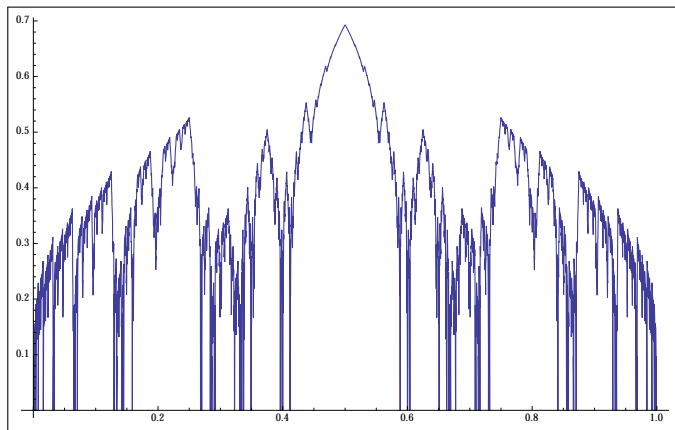
Question (Thurston, Hubbard):
Is $h(\theta)$ a continuous function of θ ?



The Main Theorem: Continuity

Theorem (T.)

The core entropy function $h(\theta)$ extends to a continuous function from \mathbb{R}/\mathbb{Z} to \mathbb{R} .



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Denote $c_j := f^j(0)$ the j^{th} iterate of the critical point, and let

$$P := \{(c_i, c_j) \mid i, j \geq 0\}$$

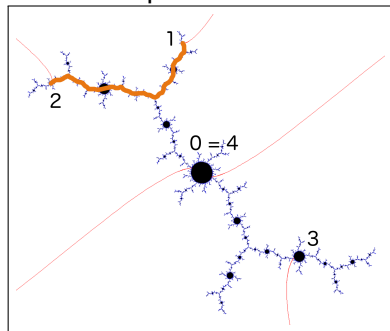
the set of pairs of postcritical points

Computing the entropy: non-separated pair

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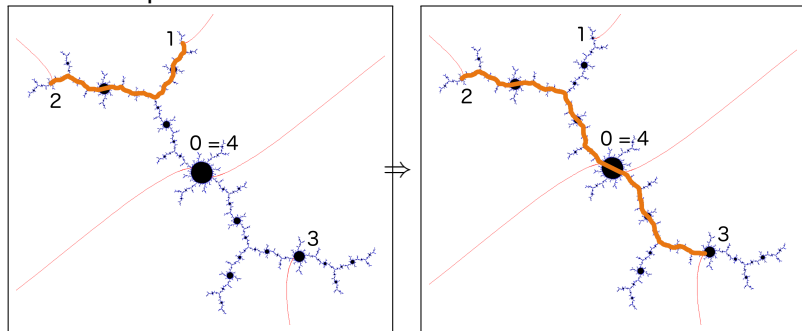
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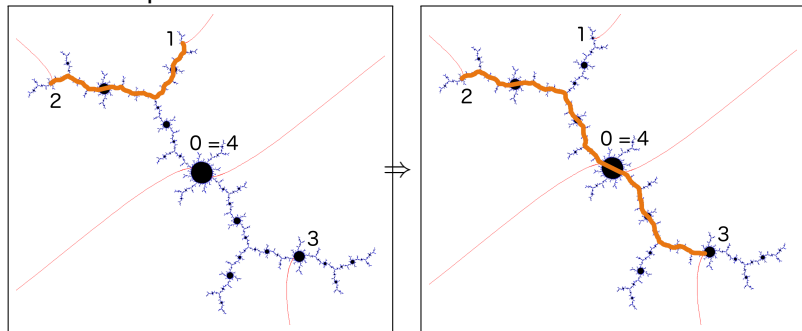
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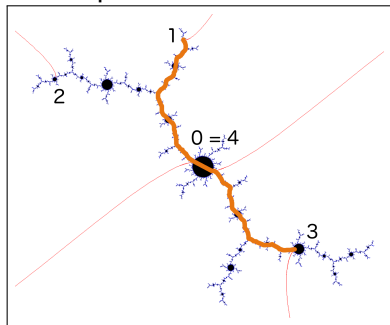
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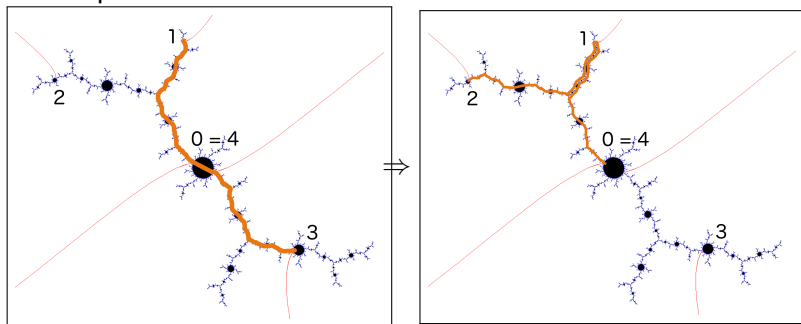
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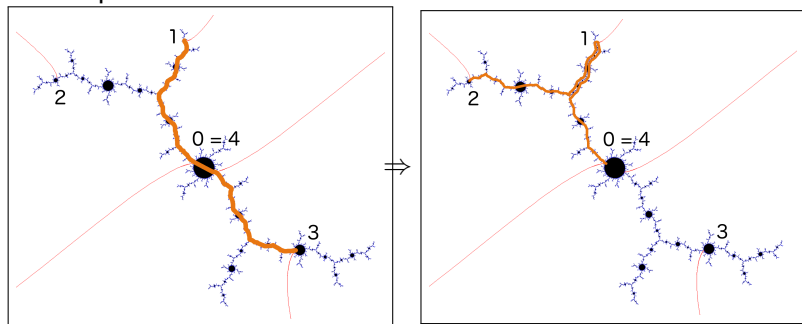
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Theorem (Thurston; Tan Lei)

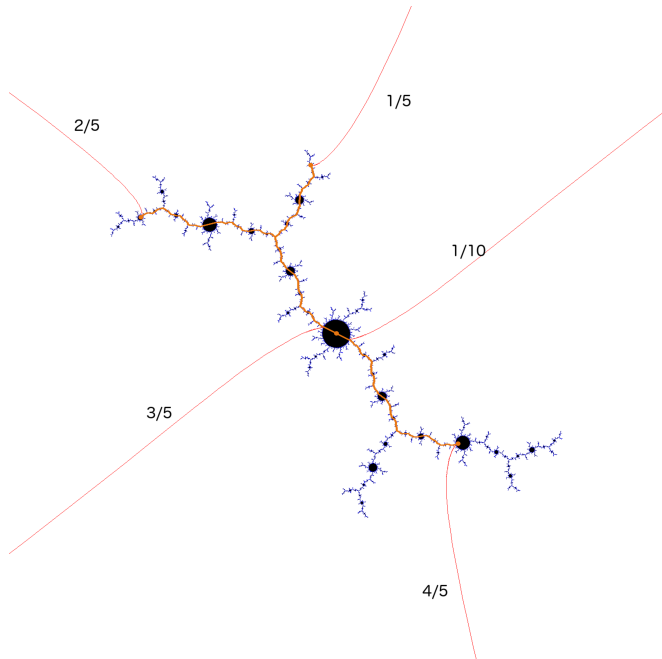
The entropy of f_θ is given by

$$h(\theta) = \log \lambda$$

where λ is the leading eigenvalue of A .

See also Gao, Jung.

The algorithm



Computing entropy: the clique polynomial

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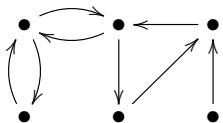
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The clique polynomial: example

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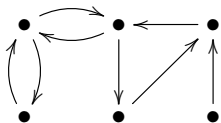
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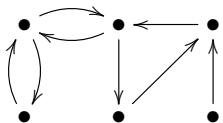
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► two 2-cycles



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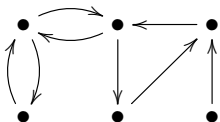
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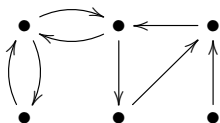
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Then we define the **growth rate** of Γ as :

$$r(\Gamma) := \limsup \sqrt[n]{C(\Gamma, n)}$$

where $C(\Gamma, n)$ is the number of closed paths of length n .

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Theorem

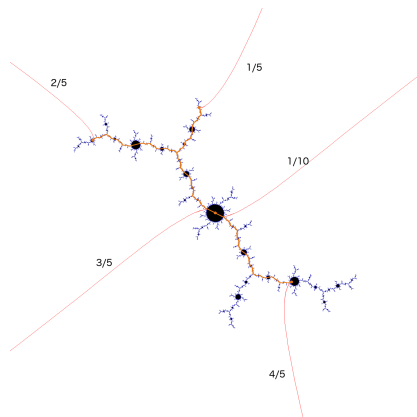
Let $\sigma \leq 1$. Then $P(t)$ defines a holomorphic function in the unit disk, and its root of minimum modulus is r^{-1} .

How to compute the core entropy without knowing complex dynamics

Let $\theta \in \mathbb{R}/\mathbb{Z}$, and $\theta_i := 2^{i-1}\theta \pmod{1}$, and consider the diameter $\{\theta/2, (\theta + 1)/2\}$ (= major leaf).

How to compute the core entropy without knowing complex dynamics

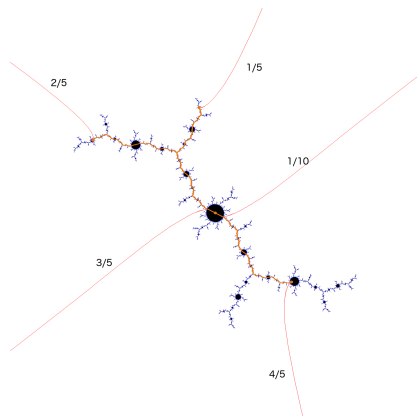
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Wedges

				...
			(4, 5)	...
		(3, 4)	(3, 5)	...
	(2, 3)	(2, 4)	(2, 5)	...
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Labeled wedges

Label all pairs as either separated or non-separated

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(The boxed pairs are the separated ones.)

From wedges to graphs

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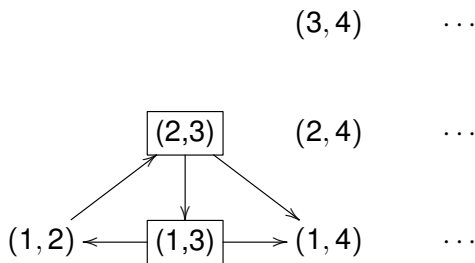
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Regularity properties of the core entropy

In fact:

Theorem (T. '15)

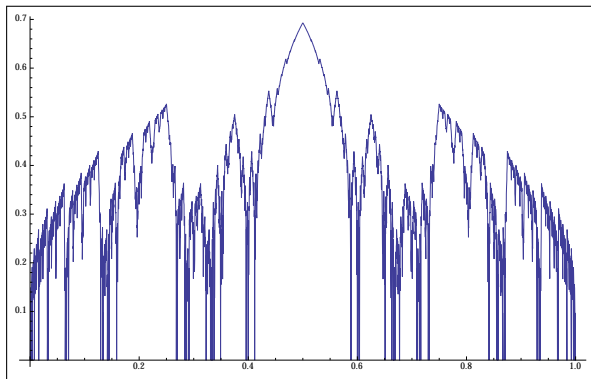
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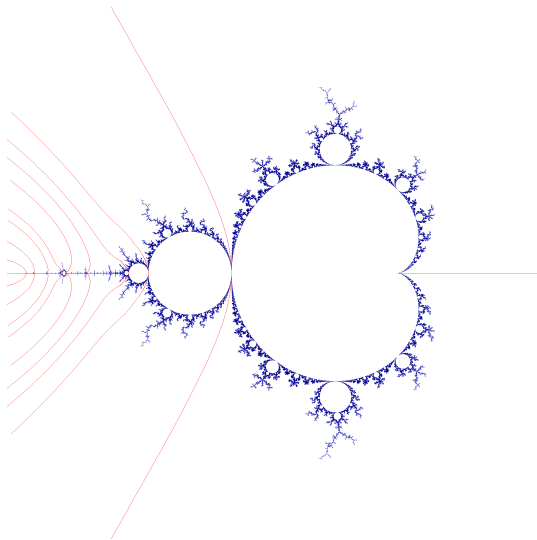
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(Conjectured Isola-Politi, 1990)

Rays landing on the real slice of the Mandelbrot set



Harmonic measure

Given a subset A of $\partial\mathcal{M}$, the **harmonic measure** $\nu_{\mathcal{M}}$ is the probability that a random ray lands on A :

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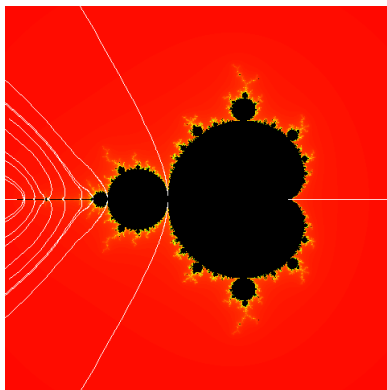
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For instance, take $A = \mathcal{M} \cap \mathbb{R}$ the real section of the Mandelbrot set. How common is it for a ray to land on the real axis?



Real section of the Mandelbrot set

Theorem (Zakeri, 2000)

*The harmonic **measure** of the real axis is 0.*

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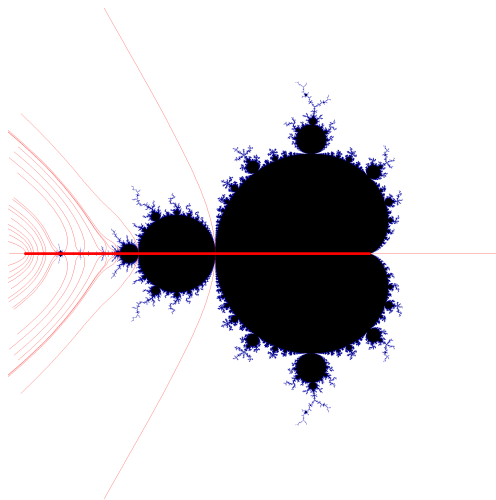
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Sectioning \mathcal{M}

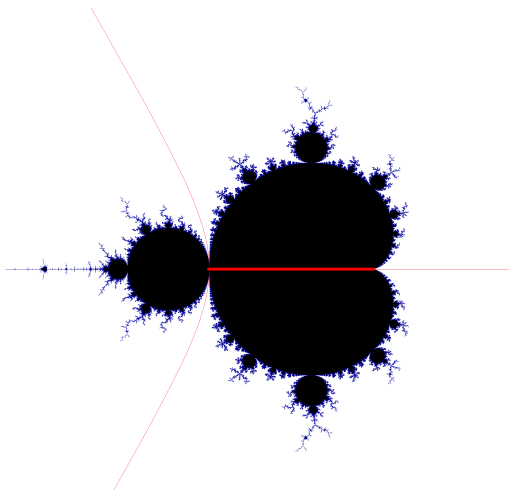
Given $c \in [-2, 1/4]$, we can consider the set of external rays which land on the real axis to the right of c :

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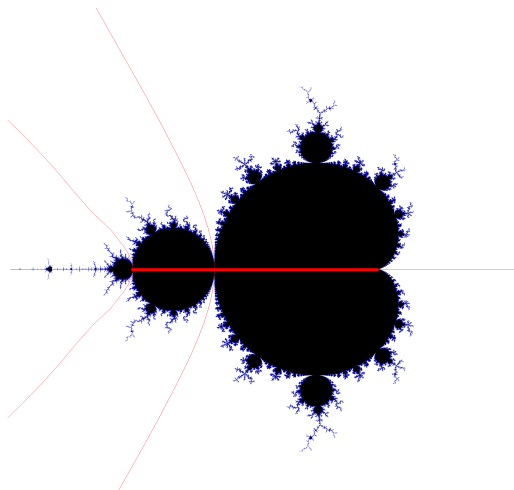
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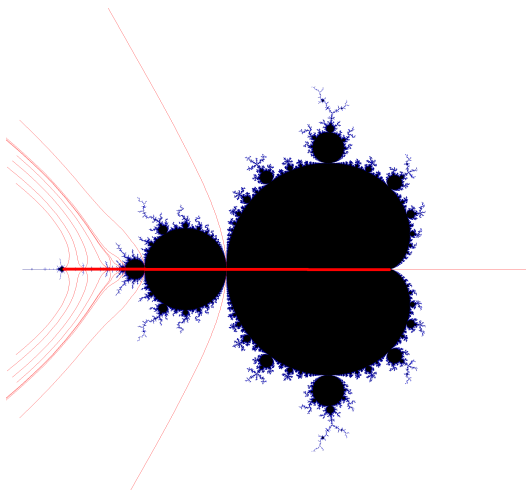
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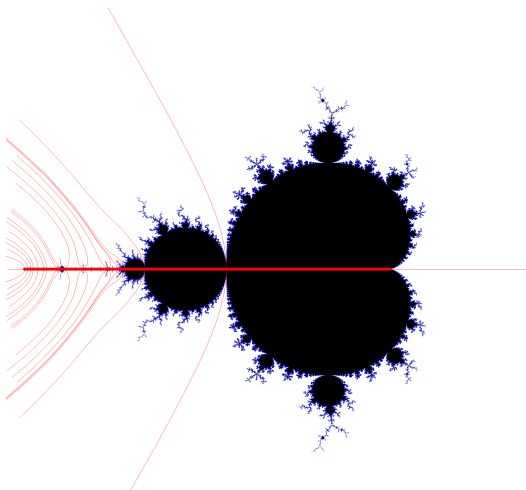
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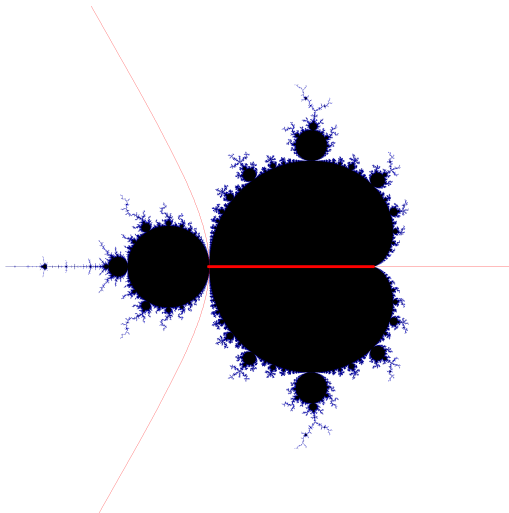


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The function

$$c \mapsto \text{H.dim } P_c$$

decreases with c , taking values between 0 and 1.

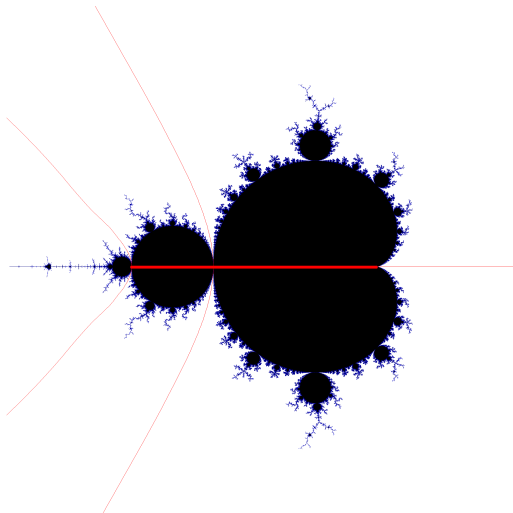


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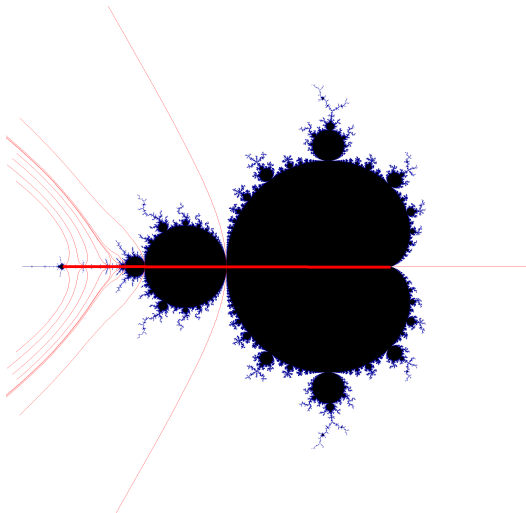


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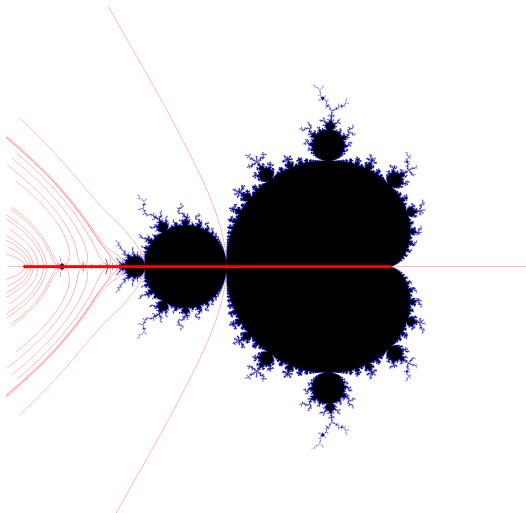


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Entropy formula, real case

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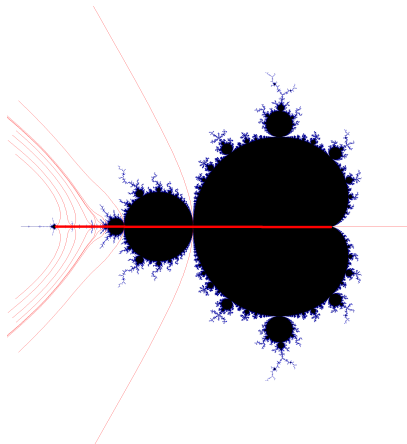
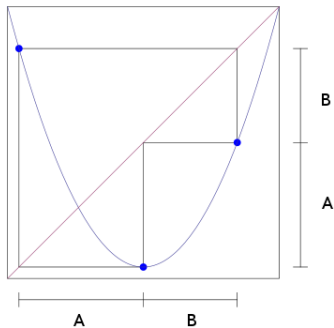
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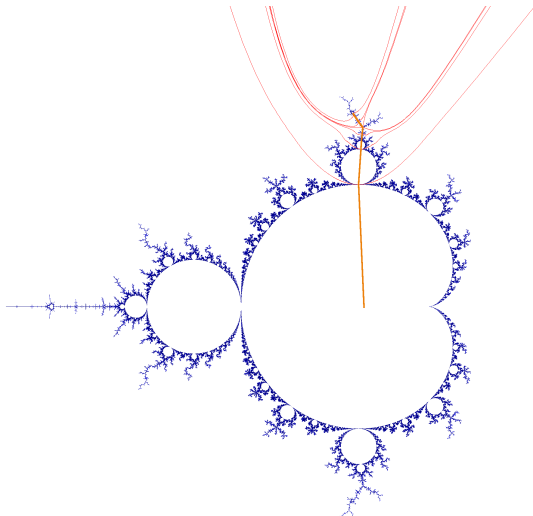
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- ▶ It can be generalized to non-real veins.

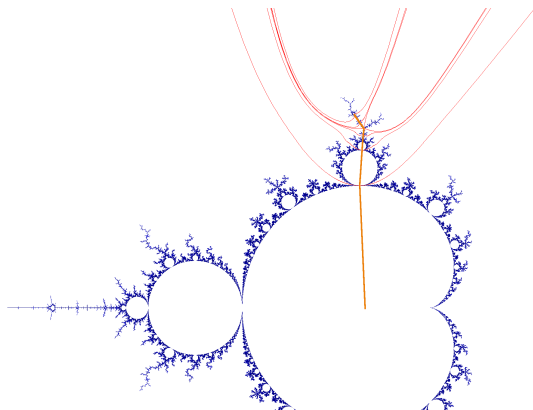
Entropy formula along complex veins

A vein is an embedded arc in the Mandelbrot set.



Entropy formula, complex case

A vein is an embedded arc in the Mandelbrot set.



Given a parameter c along a vein, we can look at the set P_c of parameter rays which land on the vein between 0 and c .

Entropy formula along complex veins

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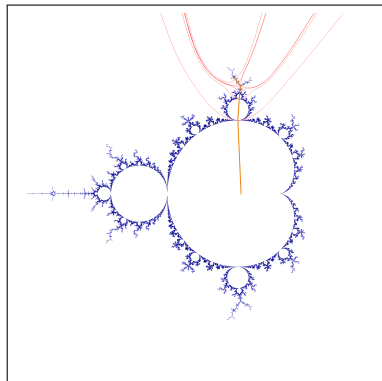
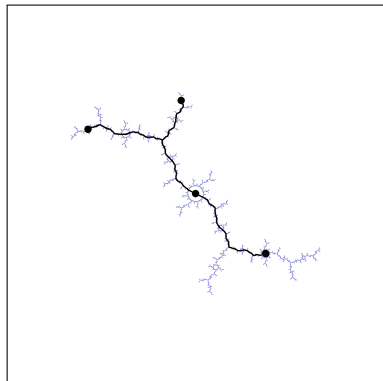
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Entropy formula along complex veins

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Further directions / questions

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4. Self-similarity of entropy function at Misiurewicz points

Laminations

Let $\theta \in \mathbb{R}/\mathbb{Z}$.

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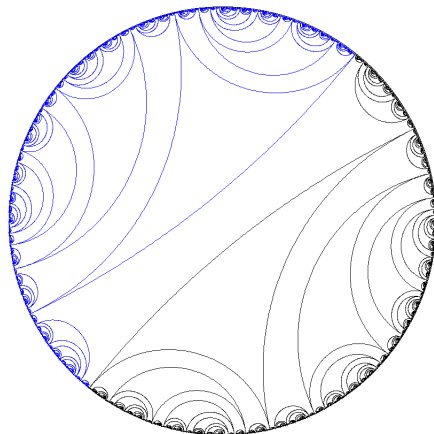
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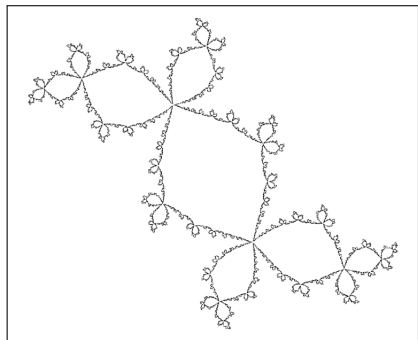
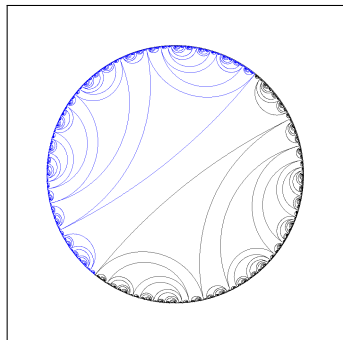
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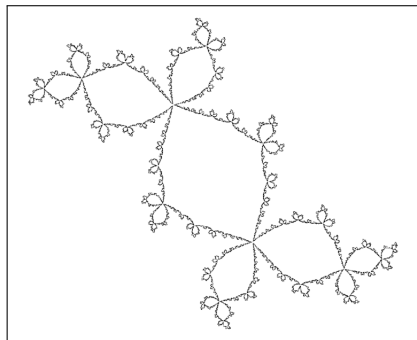
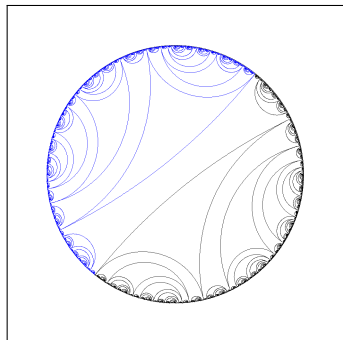
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Question. Can we define a transverse measure on \mathcal{L}_θ ?

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Let $\theta \in \mathbb{R}/\mathbb{Z}$, and $f_\theta(z) = z^2 + c_\theta$. Then there exists a lamination \mathcal{L}_θ on the disk such that $\theta_1 \sim \theta_2$ if $R(\theta_1)$ and $R(\theta_2)$ land at the same point.

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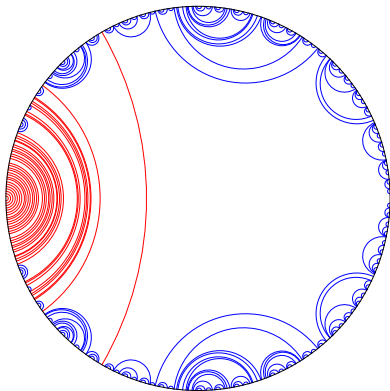
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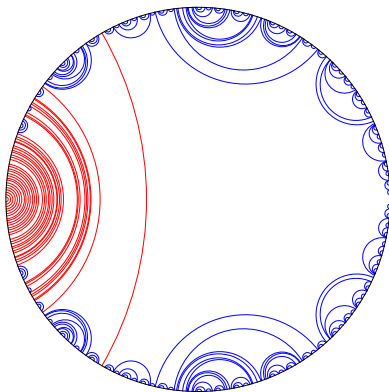
Such a measure induces a semiconjugacy between $f_\theta : T_\theta \rightarrow T_\theta$ and a piecewise linear model with slope λ_θ .
(Compare: Milnor-Thurston, Sullivan dictionary)

Thurston's quadratic minor lamination (QML)



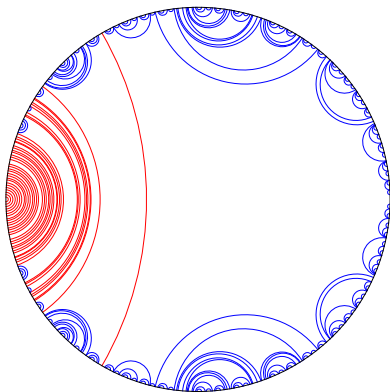
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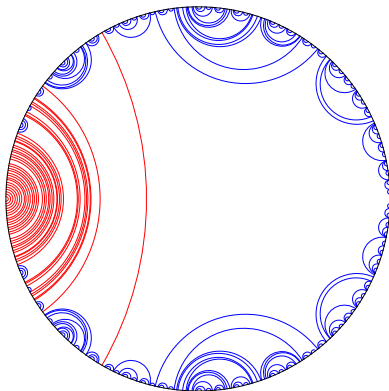
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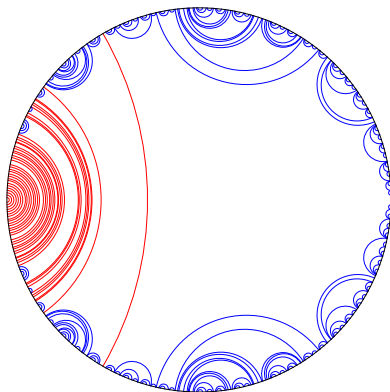
For each f_c , pick the **minor leaf** of the lamination for f_c (i.e, the ray pair landing at the critical value (or its root)). The **QML** is the union of all minor leaves for all $c \in \mathcal{M}$. The quotient \mathcal{M}_{abs} of the disk by the lamination is a (locally connected) model for the Mandelbrot set, and homeomorphic to it if MLC holds.

A transverse measure on QML



Let $l_1 < l_2$ two leaves, and τ a transverse arc connecting them.

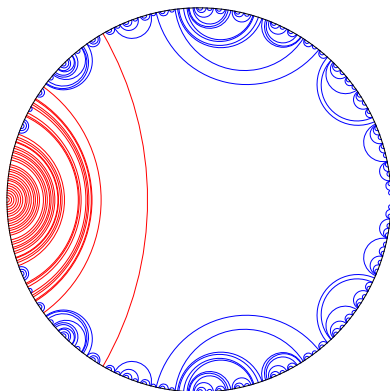
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Then we define

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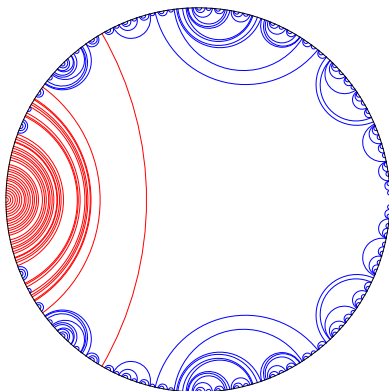


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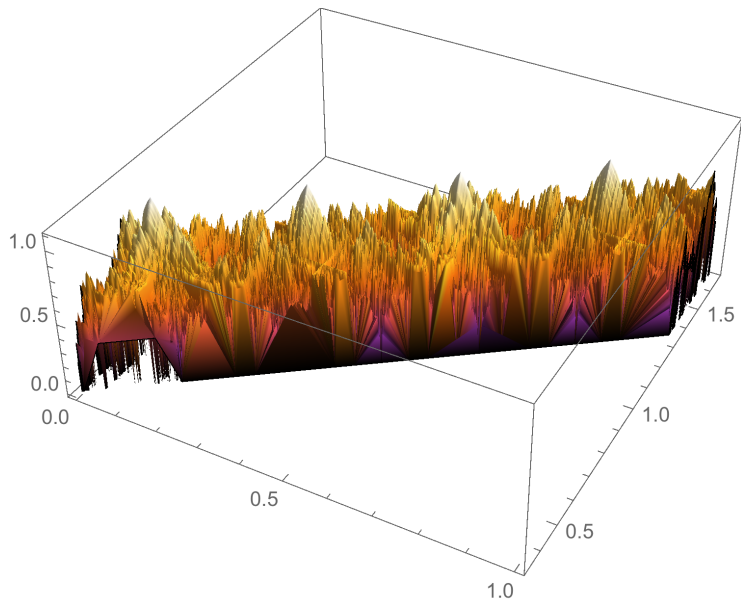
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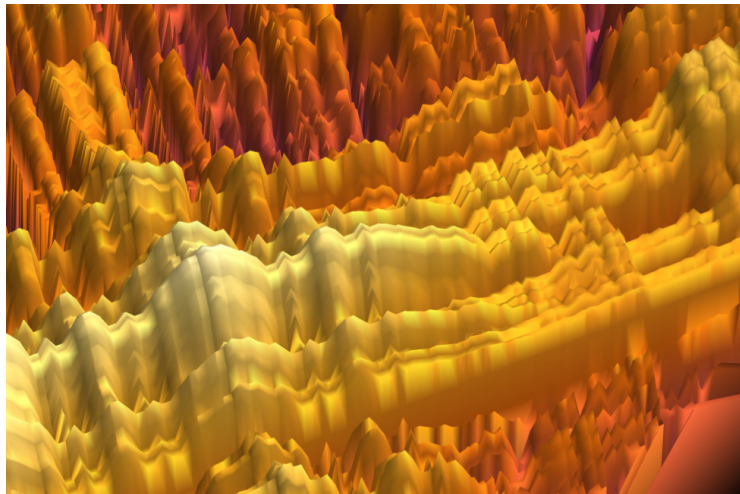
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“Combinatorial bifurcation measure”?

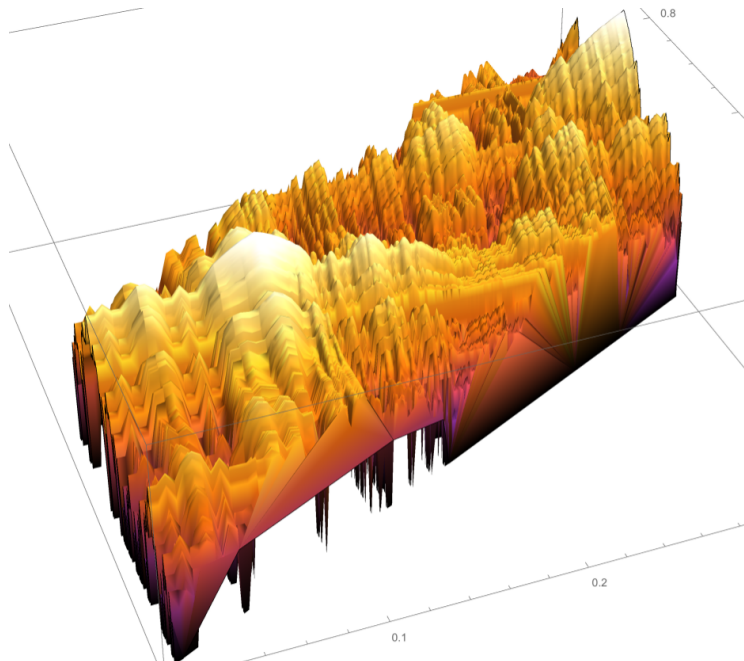
The core entropy for cubic polynomials



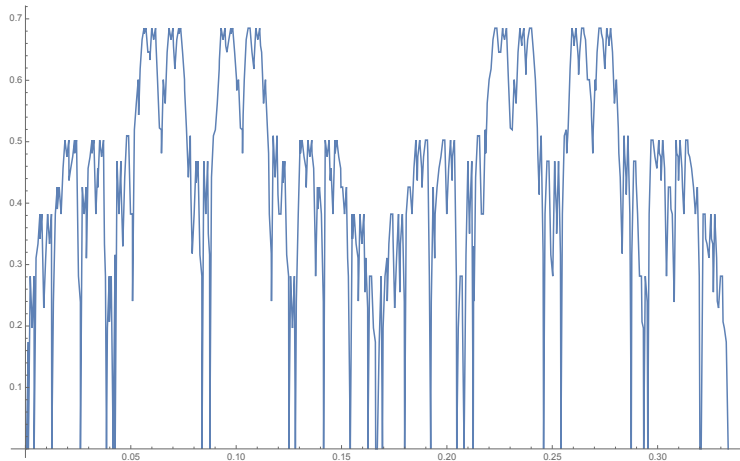
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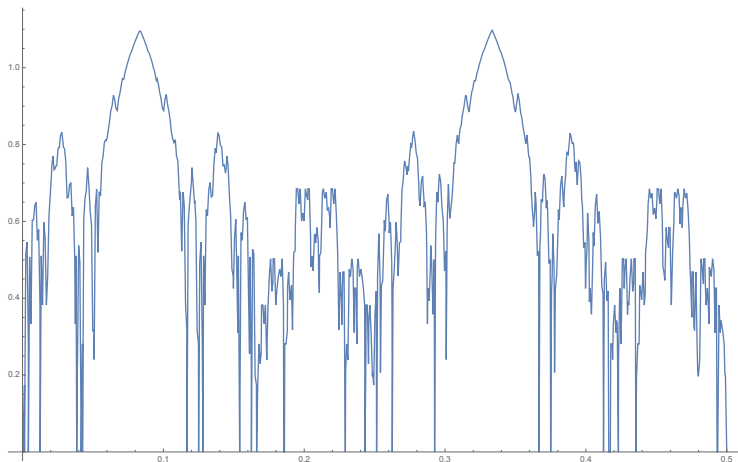


The unicritical slice



$$f(z) = z^3 + c$$

The symmetric slice



$$f(z) = z^3 + cz$$

Continuity in higher degree, combinatorial version

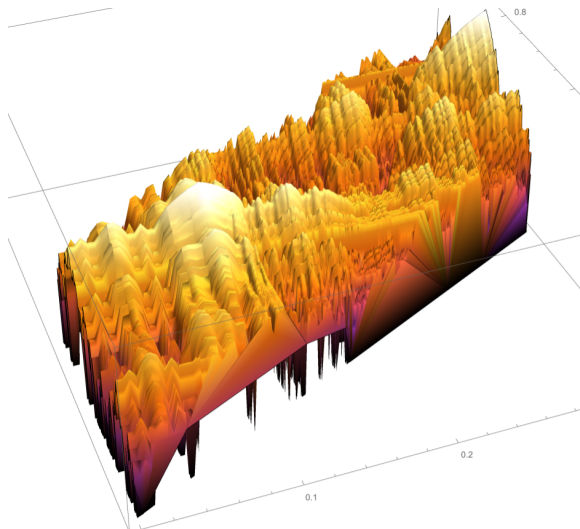
Theorem (T. - Yan Gao)

Fix $d \geq 2$. Then the core entropy extends to a continuous function on the space $\text{PM}(d)$ of primitive majors.

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The end

Thank you!